

Preferential and k -Zone Parking Functions

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Overview

- Introduction
- Relationships between Parking Functions
- Computational Results
- Other Results
- Summary

Parking Functions

- Imagine n cars enter a one-way street consisting of n parking spots and a list of parking preferences e .
- Each car entering has a preferred spot.
- If that spot is empty, then the car parks.
- If the spot is taken, then there are 5 variances of parking rules that the car can follow.

Motivating Questions

- What are the relationships between different parking functions?
- Are there any connections to other combinatorial objects?

Types of Parking Rules

$\forall S \subseteq \mathbb{R}^d, \forall L \in \mathcal{L}^d, G \sim \langle Z \rangle^S =$ Checks only forward for available spots.

$\forall Q \in \mathbb{R}^d, \forall L \in \mathcal{L}^d, G \sim \langle Z \rangle^S =$ Checks W spots backwards one at a time.

$\forall S \subseteq \mathbb{R}^d, \forall L \in \mathcal{L}^d, G \sim \langle Z \rangle^S =$ Checks back immediately W spots for availability.

$\forall Q \in \mathbb{R}^d, \forall L \in \mathcal{L}^d, G \sim \langle Z \rangle^S =$ Checks \wedge S spots back for an available spot.

$\forall f \subseteq \mathbb{R}^d, \forall Q \in \mathbb{R}^d, \forall L \in \mathcal{L}^d, G \sim \langle Z \rangle^S =$ Checks $S - 1$ spots back for an available spot.

Types of Parking Rules

$Y_{SS} - Y_d - q_{VSL} G \sim \wedge \langle z \rangle S^{\wedge} s =$ Checks only forward for available spots.

$WQ - e_{CS} d - q_{VSL} G \sim \wedge \langle z \rangle S^{\wedge} s =$ Checks W spots backwards one at a time.

$W\check{S} b^{\wedge} C d - q_{VSL} G \sim \wedge \langle z \rangle S^{\wedge} s =$ Checks back immediately W spots for availability.

$d q_{CHC}^{\wedge} z S Y_d - q_{VSL} G \sim \wedge \langle z \rangle S^{\wedge} s =$ Checks \wedge S spots back for an available spot.

$R^{\wedge} f C o p C d q_{CHC}^{\wedge} z S Y_d - q_{VSL} G \sim \wedge \langle z \rangle S^{\wedge} s =$ Checks $S - 1$ spots back for an available spot.

Classical Parking Functions

? $C^{\wedge} S^{\wedge}$

- Each car has a preferred spot which it goes to when entering the street.
- If parking spot is empty, car parks.
- Otherwise it continues down the street until it finds an empty spot to park in.

Consider the following parking preference vector $e = (2; 3; 1; 4)$:

i	p_i	Configuration
1	2	— <u>c_1</u> — —
2	3	— <u>c_1</u> <u>c_2</u> —
3	1	<u>c_3</u> <u>c_1</u> <u>c_2</u> —
4	4	<u>c_3</u> <u>c_1</u> <u>c_2</u> <u>c_4</u>

Types of Parking Rules

$Y_{SS} - Y_d - q_{WL} G \sim \wedge \langle Z \rangle S^{\wedge} S =$ Checks only forward for available spots.

$WQ - eY_{CS} d - q_{WL} G \sim \wedge \langle Z \rangle S^{\wedge} S =$ Checks W spots backwards one at a time.

$W\check{S} b^{\check{C}} d - q_{WL} G \sim \wedge \langle Z \rangle S^{\wedge} S =$ Checks back immediately W spots for availability.

$d q_{CHC}^{\wedge} z S Y_d - q_{WL} G \sim \wedge \langle Z \rangle S^{\wedge} S =$ Checks \wedge S spots back for an available spot.

$R^{\wedge} f C_{ps} C d q_{CHC}^{\wedge} z S Y_d - q_{WL} G \sim \wedge \langle Z \rangle S^{\wedge} S =$ Checks $S - 1$ spots back for an available spot.

k -Naples Parking Functions

? $C^{\wedge} S^{\$} \wedge$

- Each car prefers a spot (e_s), in which it attempts to park in.
- If spot empty, car parks.
- If occupied, car backs up checking W spots behind it's preferred spot one at a time and parks in first available.
- If there are no empty spots between e_s W and e_s then car continues down the street and parks in first available.

Consider the following parking preference vector $e = (4; 4; 3; 2; 4)$:

i	p_i	k	Configuration				
1	4	2	—	—	—	c_1	—
2	4	2	—	—	c_2	c_1	—
3	3	2	—	c_3	c_2	c_1	—
4	2	2	c_4	c_3	c_2	c_1	—
5	4	2	c_4	c_3	c_2	c_1	c_5

Types of Parking Rules

$\forall s \in S, \exists d \in D, \forall L \in L, \exists S \in S$ Checks only forward for available spots.

$\forall d \in D, \exists C \in C, \forall L \in L, \exists S \in S$ Checks W spots backwards one at a time.

$\forall S \in S, \exists C \in C, \forall L \in L, \exists S \in S$ Checks back immediately W spots for availability.

$\exists C \in C, \forall S \in S, \exists d \in D, \forall L \in L, \exists S \in S$ Checks \wedge S spots back for an available spot.

$\forall C \in C, \exists d \in D, \forall C \in C, \forall S \in S, \exists d \in D, \forall L \in L, \exists S \in S$ Checks $S - 1$ spots back for an available spot.

k-Zone Parking Functions

? $C^{\wedge} \mathbb{S} \mathbb{S}^{\wedge}$

- Each car prefers a spot (e_S), in which it attempts to park in.
- If spot is empty, car parks.
- If occupied, car backs up immediately W spots behind its preferred spot and then moves down the street if spot e_S W s taken.
- Parks in first available spot.

Consider the following parking preference vector $e = (4; 4; 3; 2; 4)$:

i	p_i	k	Configuration
1	4	2	— — — c_1 —
2	4	2	— c_2 — c_1 —
3	3	2	— c_2 c_3 c_1 —
4	2	2	c_4 c_2 c_3 c_1 —
5	4	2	c_4 c_2 c_3 c_1 c_5

k -Zone vs. k -Naples

$b^{\wedge} \lfloor z - q \rfloor$

$W - e \lfloor s - s - 4s \rfloor b^{\wedge} \lfloor s - b^{\wedge} \rfloor$

For $W=2$:

\wedge	W Naples	W Zone
1	1	1
2	4	4
3	27	27
4	240	244
5	2,731	2,808
6	38,034	39,416
7	627,405	654,302

Non-Increasing Preference Vectors

Consider the following table of the total number of parking functions formulated from non-increasing preference vector of length \hat{n} .

For $W=2$:

\hat{n}	W_{Naples}	W_{Zone}
1	1	1
2	3	3
3	10	10
4	34	34
5	117	117
6	407	407
7	1,430	1,430

Enumerating k -Zone

Consider the table enumerating the number of W Zone parking functions of length n and fixed parameter W

n	$W=0$	$W=1$	$W=2$	$W=3$	$W=4$	$W=5$	$W=6$	$W=7$
1	c							
2	3	J						
3	16	24	u					
4	125	203	244	lv				
5	1,296	2,225	2,808	3,065	{x l			
6	16,807	30,067	39,416	44,424	46,296	Jv>l v		
7	262,144	484,071	654,302	757,919	805,543	821,023	D {>J{	
8	4,782,969	9,057,316	12,553,351	14,880,368	16,110,376	16,613,896	16,757,056	cv>ll> cv

Findings From the Table (Continued)

; $b^{\wedge} \cup \{z, \dots\}$

$RH^{\wedge} 2 - \wedge @ 0 W \wedge 1 > zPC^{\wedge} j \check{S}dG(\wedge; W 1)j j \check{S}dG(\wedge; W 2)j$
 $S \cap \sim - Yzb =$

- $yPC \text{ bq@cqbHzPC} - \check{V}Cq^{\wedge} - zS^{\wedge}L Lq \sim e, \wedge_{+1}$
- $J \sim \setminus 4CqbHO - \setminus Szb^{\wedge}S^{\wedge} \langle \%dCs b^{\wedge} zPC \langle b \setminus eYzCLq eP \rangle V^{\wedge}$
- $J \sim \setminus 4CqbH^{\wedge}C \langle W \langle Cs \dots SP^{\wedge} @SzS^{\wedge} \langle z 4G @s Hbq^{\wedge}! 4G @$
 $eCq \setminus \sim z - zS^{\wedge}si$

Formula:

$$\frac{\wedge!}{2} \tag{1}$$

1; 3; 12; 60; 360; 2520; 20160; 181440; ...

coinciding with the OEIS sequence [A001710](#).

Types of Parking Rules

$Y_{SS} - Y_d - q_{VL} G \sim \wedge \langle z \rangle S^{\wedge} s =$ Checks only forward for available spots.

$W_{VQ} - e_{CS} d - q_{VL} G \sim \wedge \langle z \rangle S^{\wedge} s =$ Checks W spots backwards one at a time.

$W_{Sb}^{\check{C}} d - q_{VL} G \sim \wedge \langle z \rangle S^{\wedge} s =$ Checks back immediately W spots for availability.

$d q_{CC}^{\wedge z} S Y_d - q_{VL} G \sim \wedge \langle z \rangle S^{\wedge} s =$ Checks \wedge S spots back for an available spot.

$R^{\wedge} C_{ps} C d q_{CC}^{\wedge z} S Y_d - q_{VL} G \sim \wedge \langle z \rangle S^{\wedge} s =$ Checks $S - 1$ spots back for an available spot.

Preferential Parking Functions

? $C' \wedge \mathbb{S} \mathbb{S} \wedge$

- Each car prefers a spot (e_S), in which it attempts to park in.
- If spot is empty, car parks.
- Otherwise car checks \wedge Spots behind e_S one by one.
- If all the \wedge Spots preceding e_S are taken, car continues down the street until it finds an available spot to park in.

Consider the following parking preference vector $e = (6;6;4;3;3;3)$:

i	p_i	$n - i$	Configuration
1	6	5	— — — — — $\underline{c_1}$
2	6	4	— — — — $\underline{c_2}$ $\underline{c_1}$
3	4	3	— — — $\underline{c_3}$ $\underline{c_2}$ $\underline{c_1}$
4	3	2	— — $\underline{c_4}$ $\underline{c_3}$ $\underline{c_2}$ $\underline{c_1}$
5	3	1	— $\underline{c_5}$ $\underline{c_4}$ $\underline{c_3}$ $\underline{c_2}$ $\underline{c_1}$
6	3	0	— $\underline{c_5}$ $\underline{c_4}$ $\underline{c_3}$ $\underline{c_2}$ $\underline{c_1}$

Types of Parking Rules

$Y_{SS} - Y_{d-q} \wedge L \sim G \sim \wedge \langle Z \rangle S \wedge S =$ Checks only forward for available spots.

$W \wedge - e \wedge C s \wedge d - q \wedge L \sim G \sim \wedge \langle Z \rangle S \wedge S =$ Checks W spots backwards one at a time.

$W \wedge b \wedge C \wedge d - q \wedge L \sim G \sim \wedge \langle Z \rangle S \wedge S =$ Checks back immediately W spots for availability.

$d \wedge C \wedge C \wedge \wedge Z \wedge S \wedge Y_{d-q} \wedge L \sim G \sim \wedge \langle Z \rangle S \wedge S =$ Checks \wedge S spots back for an available spot.

$R \wedge f \wedge C \wedge C \wedge d \wedge C \wedge C \wedge \wedge Z \wedge S \wedge Y_{d-q} \wedge L \sim G \sim \wedge \langle Z \rangle S \wedge S =$ Checks $S - 1$ spots back for an available spot.

Inverse Preferential Parking Functions

? $C' \wedge S \mathbb{S} \wedge$

- Each car prefers a spot (e_S), in which it attempts to park in.
- If spot is empty, car parks.
- If occupied, car checks $S - 1$ spots behind e_S one by one.
- If all $S - 1$ spots behind e_S are taken, car continues down the street until it finds an available spot to park in.

Consider the following parking preference vector $e = (6; 6; 4; 3; 3; 3)$:

i	p_i	$i - 1$	Configuration
1	6	0	— — — — — c_1
2	6	1	— — — — — c_2 c_1
3	4	2	— — — c_3 c_2 c_1
4	3	3	— — — c_4 c_3 c_2 c_1
5	3	4	— c_5 c_4 c_3 c_2 c_1
6	3	5	c_6 c_5 c_4 c_3 c_2 c_1

Preferential Parking Function Findings

$X \setminus -$

$R^H \quad 2 > zPC^{\wedge} (\wedge; \dots; \wedge) \quad 2 [\wedge]^{\wedge} \quad S^{\wedge} bz - eq \quad C^{\wedge} zS \quad Ye - q \quad B^{\wedge} L \quad H^{\wedge} zS^{\wedge} i$

$d \phi b H$

By contradiction assuming that all cars can park.

Other Results

$X \setminus -$

, $Y \text{ eq } H \text{ eq } C^{\wedge} < C \text{ f } C < z \text{ b } q \text{ s } - q \text{ C } S \text{ f } C \text{ o } p \text{ C } \text{ eq } H \text{ eq } C^{\wedge} z \text{ S } Y \text{ e } - q \text{ S } L \text{ H } ^{\wedge} < z \text{ S } ^{\wedge} \text{ si}$

d o p b H

Direct proof.

